

Computing supersingular isogenies on Kummer surfaces

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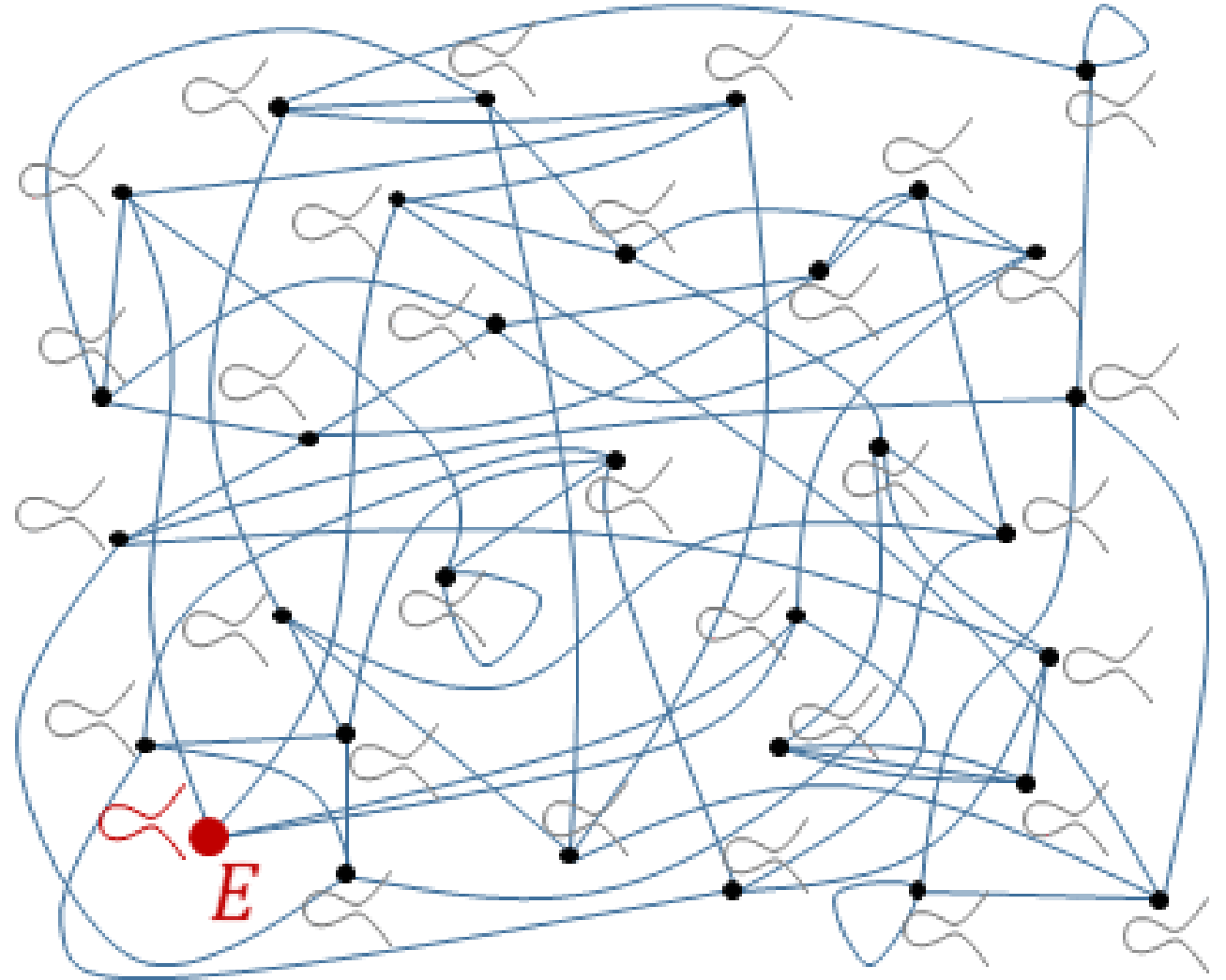
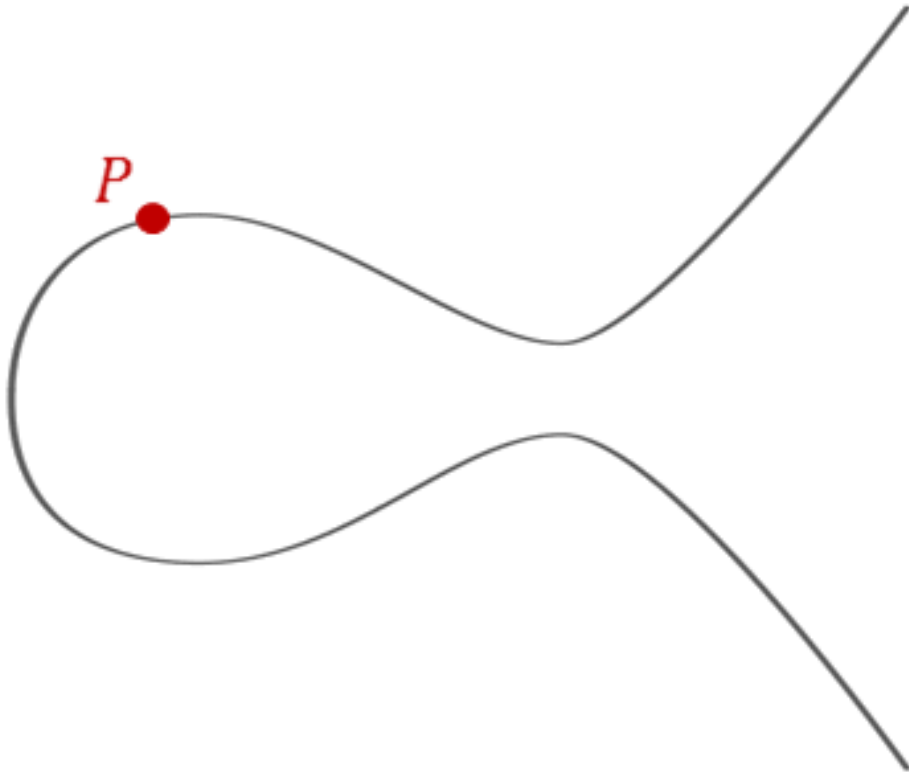


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ECC

vs.

post-quantum ECC



Alice 2^e -isogenies, Bob 3^f -isogenies



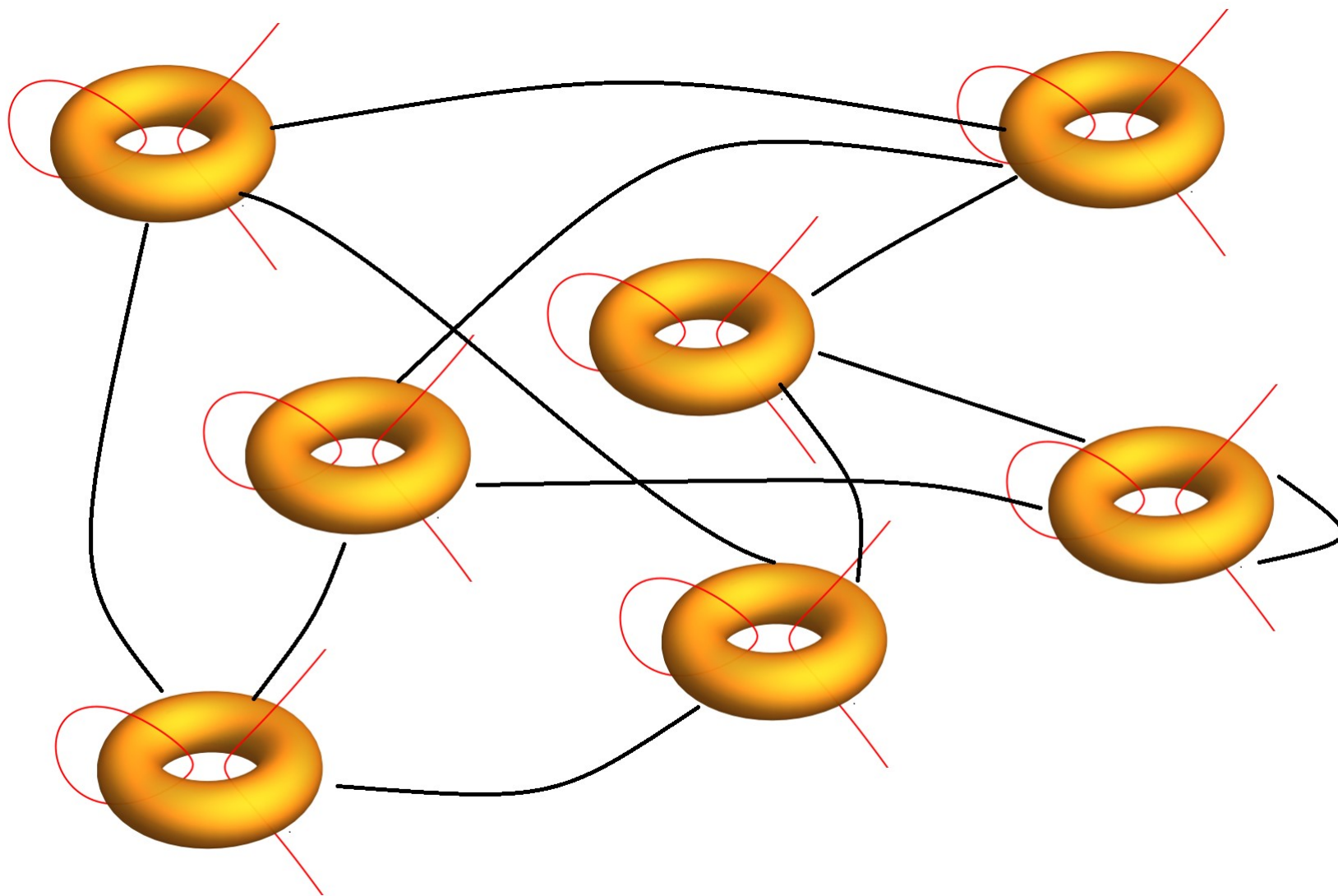
Alice



Bob

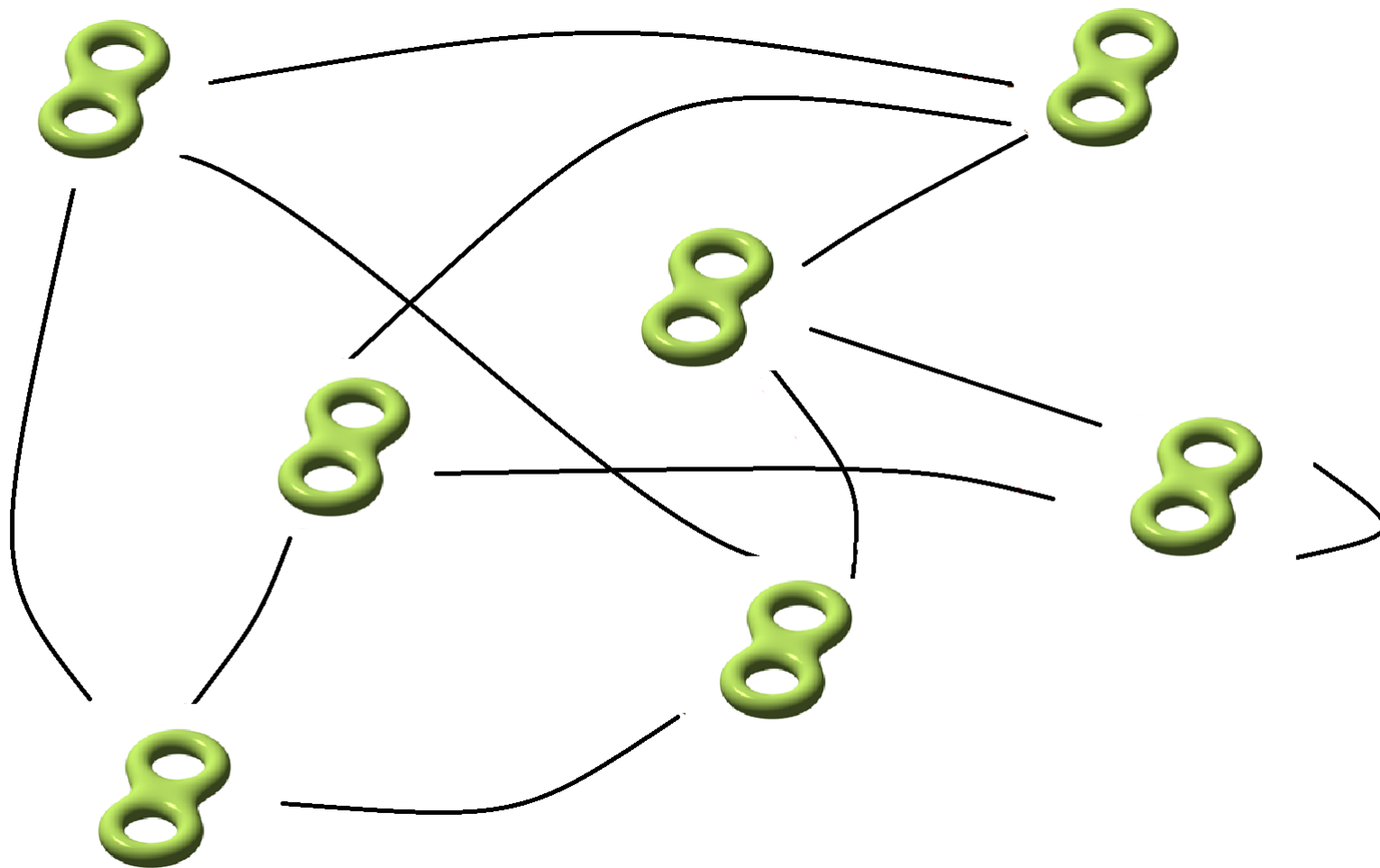
In a nutshell:

$$E(\mathbb{F}_{p^2})$$



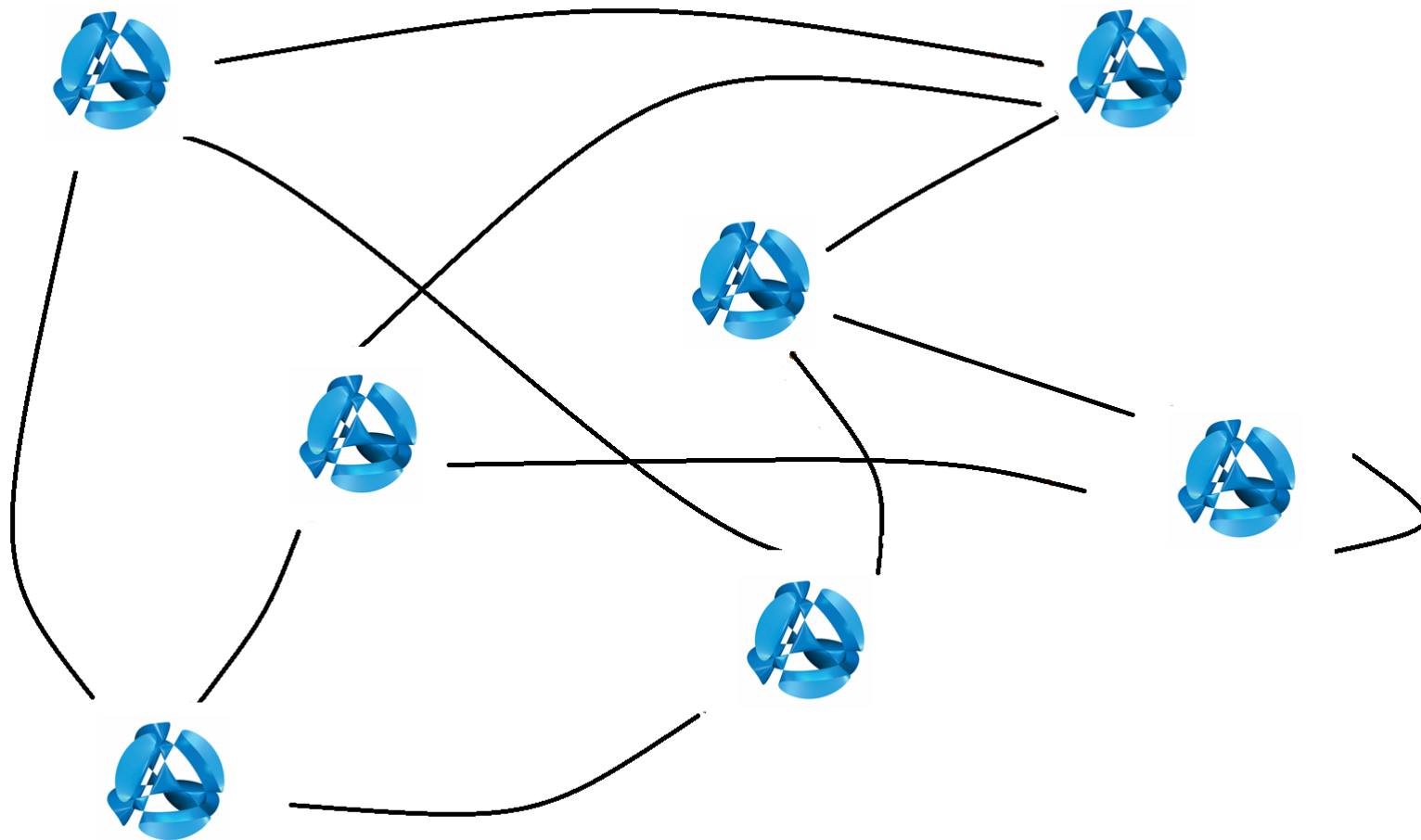
In a nutshell:

$$J_C(\mathbb{F}_p)$$

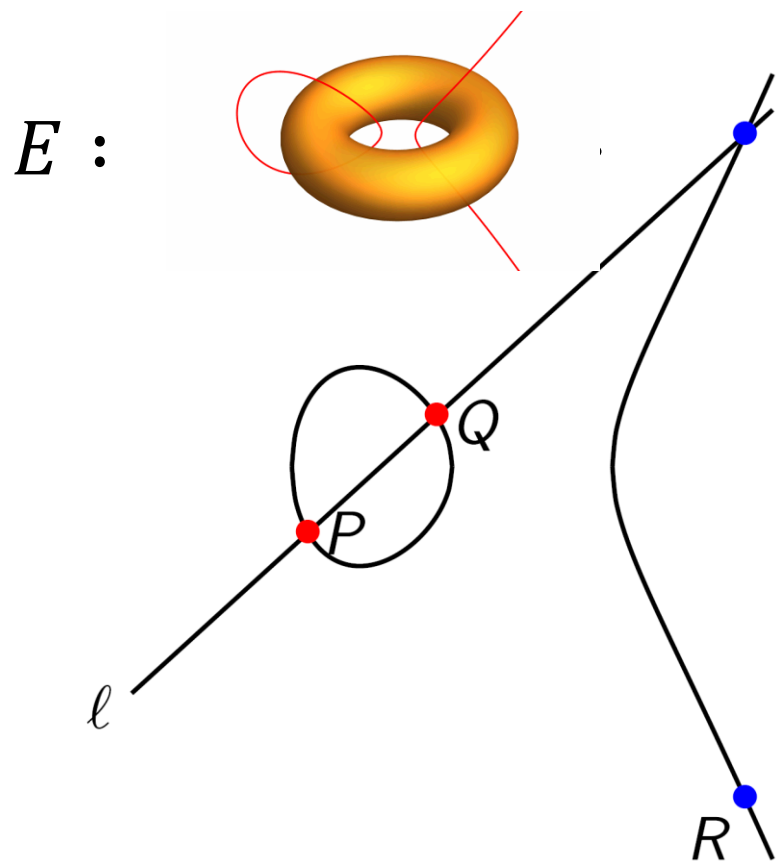


In a nutshell:

$$K(\mathbb{F}_p)$$

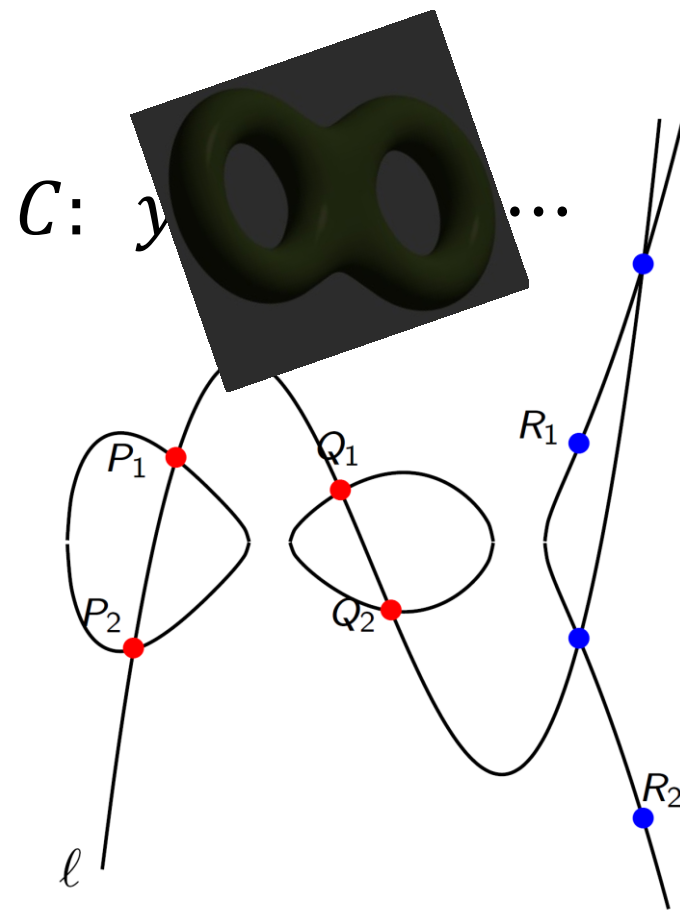


Why go hyperelliptic?



$$G = E$$
$$|G| = \#E$$

$$\#E(\mathbb{F}_q) \approx \#C(\mathbb{F}_q)$$



$$G \approx C \times C$$
$$|G| = (\#C)^2$$

Why go Kummer?



72 equations in \mathbb{P}^{15}



$K(\mathbb{F}_p) = J(\mathbb{F}_p) / \langle \pm 1 \rangle$
1 equation in \mathbb{P}^3

- Genus 2 analogue of elliptic curve x -line
- Extremely efficient arithmetic

... a few of my favourite things...

WEIL RESTRICTION OF AN ELLIPTIC CURVE OVER A QUADRATIC EXTENSION

JASPER SCHOLTEN

ABSTRACT. Let K be a finite field of characteristic not equal to 2, and L a quadratic extension of K . For a large class of elliptic curves E defined over L we construct hyperelliptic curves over K of genus 2 whose jacobian is isogenous to the Weil restriction $\text{Res}_K^L(E)$.

Hyper-and-elliptic-curve cryptography

Daniel J. Bernstein and Tanja Lange

At this point one can and should object that [48, Lemma 2.1] merely guarantees the existence of an isogeny from W to J ; it does not guarantee the existence of an *efficient* isogeny from W to J .

The main challenge addressed in this section is to show that W and J are *efficiently* isogenous.

Fast genus 2 arithmetic based on Theta functions

P. Gaudry

Remark 3.5. The pseudo-group law that we just described is somewhat surprising, because it heavily relies on a (2,2)-isogenous abelian variety for the computation: for the doubling, the point is pushed through isogenies back and forth, thus obtaining a multiplication by 2 map.

TOWARDS QUANTUM-RESISTANT CRYPTOSYSTEMS FROM SUPERSINGULAR ELLIPTIC CURVE ISOGENIES

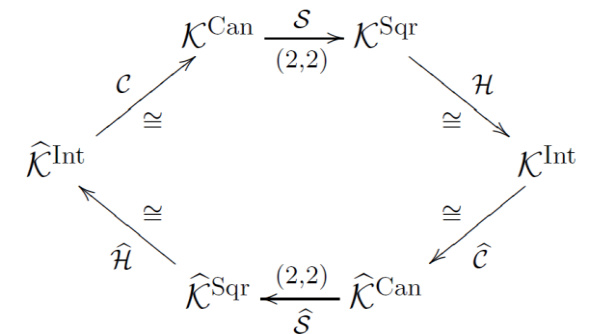
LUCA DE FEO, DAVID JAO, AND JÉRÔME PLÛT

Also observe that since P and $-P$ generate the same subgroup, isogenies can be defined and evaluated correctly on the Kummer line.

It is not immediately evident how to put F in Montgomery form without computing square roots. If P_8 is a point satisfying $[2]P_8 = P_4$, then $\phi(P_8) = (2\sqrt{2+A}, \dots)$, and F can be put in the form

qDSA: Small and Secure Digital Signatures with Curve-based Diffie–Hellman Key Pairs

Joost Renes^{1*} and Benjamin Smith²



From elliptic to hyperelliptic

Consider

$$E/K: y^2 = x^3 + 1$$

$$C/K: y^2 = x^6 + 1$$

Obvious map

$$\begin{aligned} \omega : C(K) &\rightarrow E(K) \\ (x, y) &\mapsto (x^2, y) \end{aligned}$$

- 1: But what about $\omega^{-1} : E(K) \rightarrow C(?)$...
- 2: Points on E are group elements, points on C are not...
- 3: Actually want map $E \rightarrow J_C$, but $\dim(E) = 1$ while $\dim(J_C) = 2$...
- 4: Want *general* ω, ω^{-1} between $y^2 = x^3 + Ax^2 + x$ to $y^2 = x^6 + Ax^4 + x^2$???

Proposition 1

$$\mathbb{F}_{p^2} = \mathbb{F}_p(i) \text{ with } i^2 + 1 = 0$$

$$E/\mathbb{F}_{p^2}: y^2 = x(x - \alpha)(x - 1/\alpha)$$

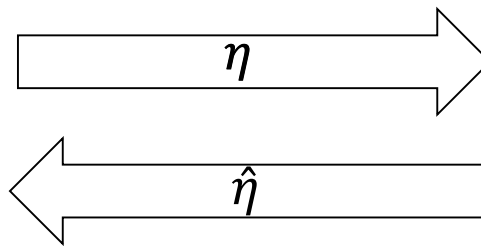
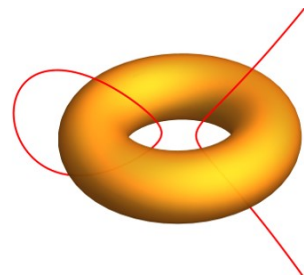
$$\alpha = \alpha_0 + \alpha_1 i \text{ with } \alpha_0, \alpha_1 \in \mathbb{F}_p$$

$$C/\mathbb{F}_p: y^2 = (x^2 + mx - 1)(x^2 - mx - 1)(x^2 - mnx - 1)$$

$$m = \frac{2\alpha_0}{\alpha_1}, n = \frac{(\alpha_0^2 + \alpha_1^2 - 1)}{(\alpha_0 + \alpha_1^2 + 1)} \text{ both in } \mathbb{F}_p$$

Then $\text{Res}_{\mathbb{F}_{p^2}/\mathbb{F}_p}(E)$ is $(2,2)$ -isogenous to $J_C(\mathbb{F}_p)$

Or, pictorially,



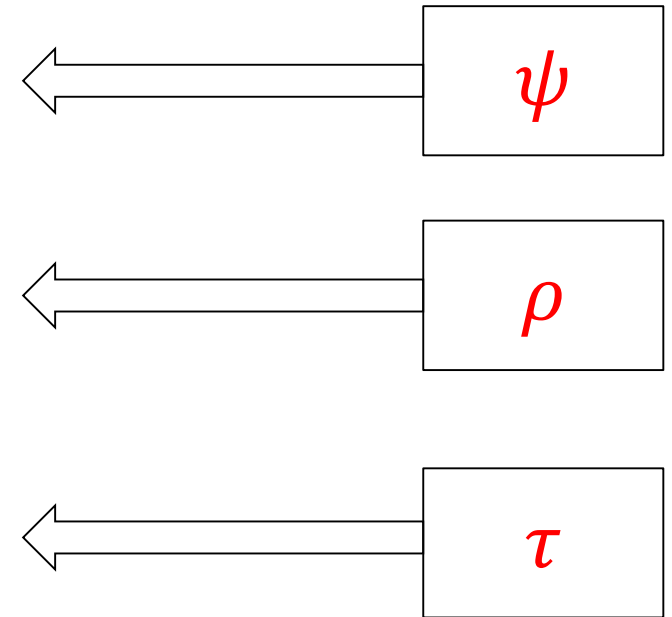
$$\ker(\eta) \cong \ker(\hat{\eta}) \cong \mathbb{Z}_2 \times \mathbb{Z}_2$$

$$\eta \circ \hat{\eta} = [2]$$

Unpacking Proposition 1

- Weil restriction turns 1 equation over \mathbb{F}_{p^2} into two equations over \mathbb{F}_p
- Simple linear transform of $E/\mathbb{F}_{p^2}: y^2 = f(x) = x^3 + Ax^2 + x$ to $\tilde{E}/\mathbb{F}_{p^2}: y^2 = g(x)$ such that $C/\mathbb{F}_{p^2}: y^2 = g(x^2)$ is non-singular
- Pullback ω^* of $\omega : (x, y) \mapsto (x^2, y)$ gives 2 points in $C(\mathbb{F}_{p^4})$, but composition with Abel-Jacobi map bring these to $J_C(\mathbb{F}_{p^2})$
- Need to go from $J_C(\mathbb{F}_{p^2})$ to $J_C(\mathbb{F}_p)$; cue good old Trace map,

$$\tau: P \mapsto \sum_{\sigma \in \text{Gal}(\mathbb{F}_{p^2}/\mathbb{F}_p)}^n \sigma(P)$$



$$\eta : \text{Res}_{\mathbb{F}_{p^2}/\mathbb{F}_p}(E) \rightarrow J_C(\mathbb{F}_p), \quad P \mapsto (\tau \circ \rho \circ \psi)(P)$$

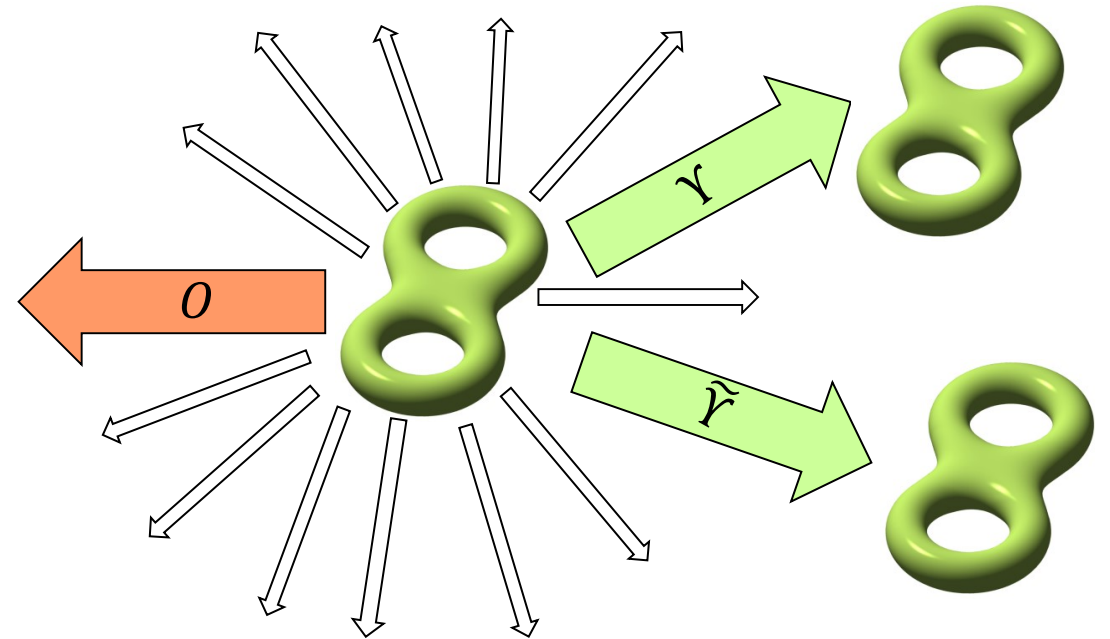
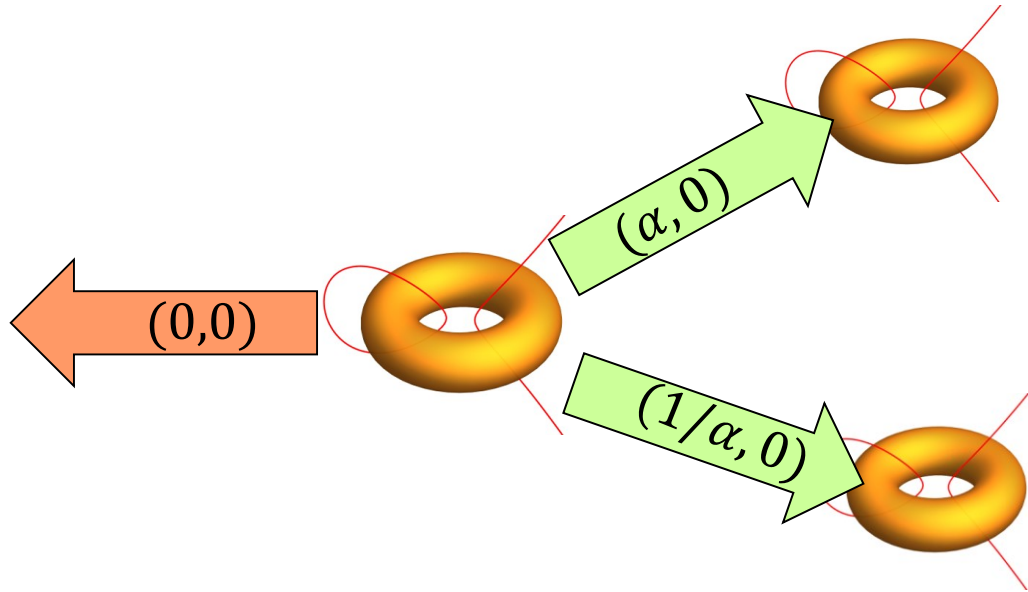
Matching 2-kernels in \mathbb{F}_{p^2} with (2,2)-kernels in \mathbb{F}_p

$$E \cong \mathbb{Z}_{(p+1)} \times \mathbb{Z}_{(p+1)}$$

$$E[2] \cong \mathbb{Z}_2 \times \mathbb{Z}_2$$

$$J_C \cong \mathbb{Z}_{(p+1)/2} \times \mathbb{Z}_{(p+1)/2} \times \mathbb{Z}_2 \times \mathbb{Z}_2$$

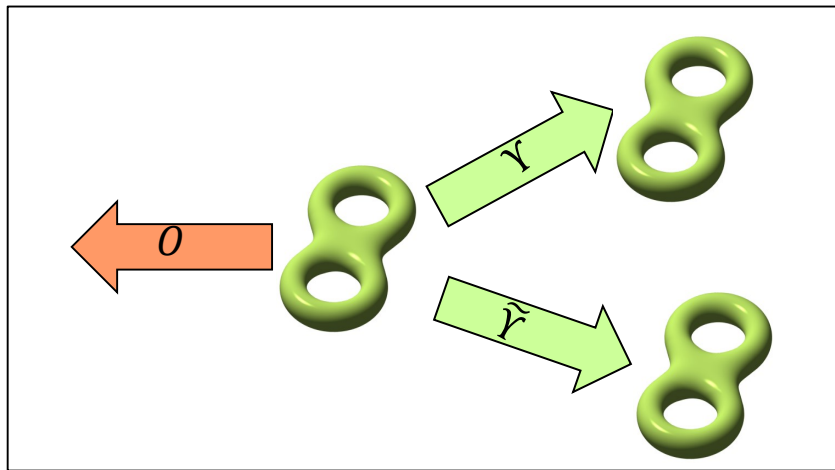
$$J_C[2] \cong \mathbb{Z}_2 \times \mathbb{Z}_2 \times \mathbb{Z}_2 \times \mathbb{Z}_2$$



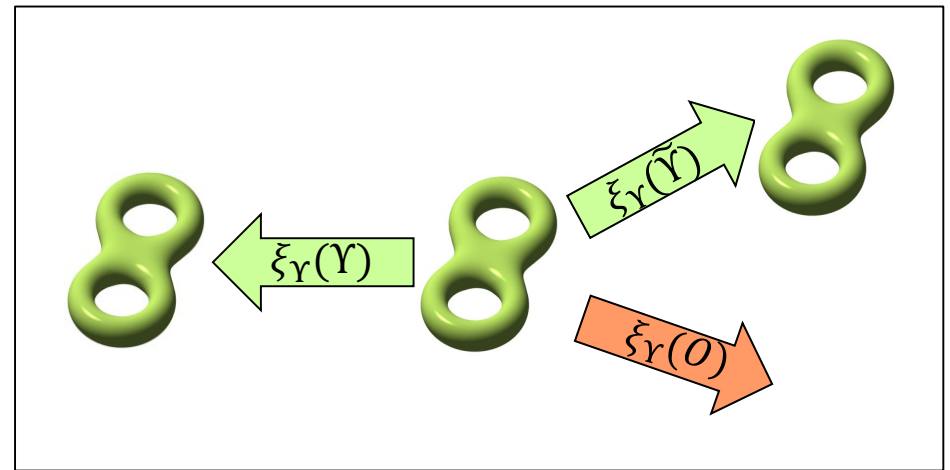
- Fifteen (2,2)-kernels in $J_C(\mathbb{F}_p)$. Number of ways to split C 's sextic into three quadratic factors.
- Lemma 2: identifies $0 \leftrightarrow (0, 0)$ and $\{\gamma, \tilde{\gamma}\} \leftrightarrow \{(\alpha, 0), (1/\alpha, 0)\}$

Richelot isogenies in genus 2

- Elliptic curve isogenies are easy/explicit/fast, thanks to Vélu. But beyond elliptic curves, far from true!
- (2,2)-isogenies in genus 2 are exception, thanks to work beginning with Richelot in 1836
- Lessons learned from elliptic case:
 - (1) easiest to derive explicitly when the kernel is \mathcal{O} , i.e. the kernel we don't want!
 - (2) when kernel is \mathcal{Y} , precompose with isomorphism $\xi_{\mathcal{Y}} : J_C \rightarrow J_{C'}$, $\mathcal{Y} \mapsto \mathcal{O}'$
 - (3) $\xi_{\mathcal{Y}}$ either requires a square root, or torsion "from above"
 - (4) who cares about the full Jacobian group, let's move the Kummer variety



$$\xi_{\mathcal{Y}} \cong$$



Supersingular Kummer surfaces

$$K_{F,G,H}^{\text{Sqr}}: F \cdot X_1 X_2 X_3 X_4 = \left(X_1^2 + X_2^2 + X_3^2 + X_4^2 - G(X_1 + X_2)(X_3 + X_4) - H(X_1 X_2 + X_3 X_4) \right)^2$$

Surface constants $F, G, H \in \mathbb{F}_p$

Points $(X_1 : X_2 : X_3 : X_4) \in \mathbb{P}^3(\mathbb{F}_p)$

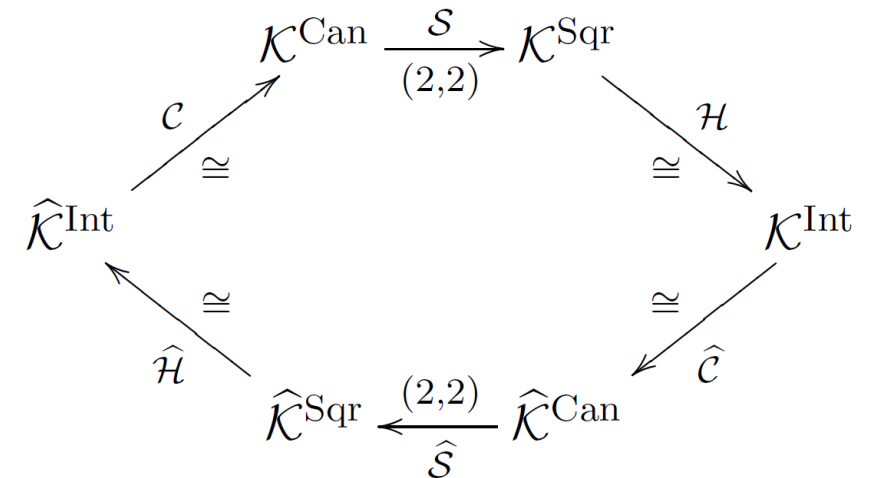
Theta constants $(\mu_1 : \mu_2 : 1 : 1) \sim (\lambda \mu_1 : \lambda \mu_2 : \lambda : \lambda)$

Arithmetic constants $(\pi_1 : \pi_2 : \pi_3 : \pi_4)$; functions of μ_1, μ_2

$$S: (\ell_1 : \ell_2 : \ell_3 : \ell_4) \mapsto (\ell_1^2 : \ell_2^2 : \ell_3^2 : \ell_4^2)$$

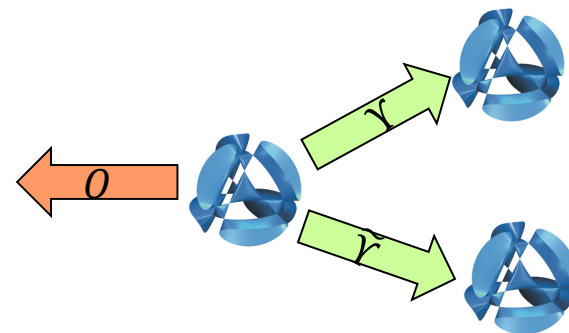
$$C: (\ell_1 : \ell_2 : \ell_3 : \ell_4) \mapsto (\pi_1 \ell_1 : \pi_2 \ell_2 : \pi_3 \ell_3 : \pi_4 \ell_4)$$

$$H: (\ell_1 : \ell_2 : \ell_3 : \ell_4) \mapsto (\ell_1 + \ell_2 + \ell_3 + \ell_4 : \ell_1 + \ell_2 - \ell_3 - \ell_4 : \ell_1 - \ell_2 + \ell_3 - \ell_4 : \ell_1 - \ell_2 - \ell_3 + \ell_4)$$



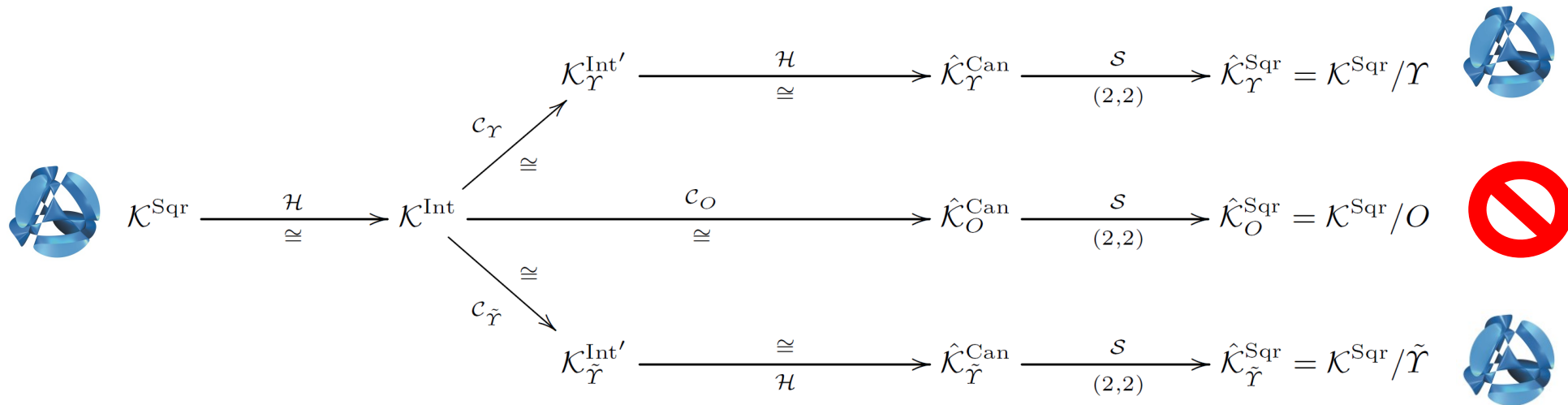
Doubling $[2]_{K^{\text{Sqr}}}: P \mapsto (S \circ \widehat{C} \circ H \circ S \circ C \circ H)(P)$

2-isogeny (splitting $[2]$) $\varphi_0: P \mapsto (S \circ C \circ H)(P)$



Kummer isogenies for non-trivial kernels

- P point of order 2 on K corresponding to $G \in \{\Upsilon, \tilde{\Upsilon}\}$. Write $H(P) = (P'_1 : P'_2 : P'_3 : P'_4)$
- Q point of order 4 on K such that $[2]Q = P$. Write $H(Q) = (Q'_1 : Q'_2 : Q'_3 : Q'_4)$
- Define $C_{Q,P} : (X_1 : X_2 : X_3 : X_4) \mapsto (\pi'_1 X_1 : \pi'_2 X_2 : \pi'_3 X_3 : \pi'_4 X_4)$
 where $(\pi_1 : \pi_2 : \pi_3 : \pi_4) = (P'_2 Q'_4 : P'_1 Q'_4 : P'_2 Q'_1 : P'_2 Q'_1)$
- Then $\varphi_P : K^{Sqr} \rightarrow K^{Sqr}/G$, $P \mapsto (S \circ H \circ C_{Q,P} \circ H)(P)$ 4M+4S+16A



Implications

Operation	chained 2-isogenies on Montgomery curves over \mathbb{F}_{p^2} (previous work)				chained (2, 2)-isogenies on Kummer surfaces over \mathbb{F}_p (this work)				
	M	S	A	\approx cycles	m	s	a	\approx cycles	
								s = m	s = 0.8m
doubling	4	2	4	5862	8	8	16	6272	5714
2-isog. curve	-	2	1	2088	19	4	28	9231	8952
2-isog. point	4	0	4	4336	4	4	16	3480	3200

- Theta constants map to theta constants: no special map needed to find image surface
- Comparison in Table/paper very conservative. Kummer will win in aggressive impl.:
 - Recall Kummer over $\mathbb{F}_{2^{127}-1}$ almost as fast as FourQ over $\mathbb{F}_{(2^{127}-1)^2}$ (scalars 4 x larger)
 - Recall that "doubling" and "2-isog. point" are bottlenecks in optimal tree strategy
 - Pushing points through 2^ℓ for small ℓ likely to be better on Kummer, don't need to compute all intermediate surface constants

Related future work

- To use this right now, Alice need to map back-and-forth using η and $\hat{\eta}$. Certainly not a deal-breaker! **Thus, this is a call for skilled implementers!**
- But ideally we want Bob to be able to use the Kummer, too! Then uncompressed SIDH/SIKE can be defined as Kummer everywhere!
Thus, this is a call for fast $(3, 3)$ -isogenies on fast Kummers!
- Going further, general isogenies in Montgomery elliptic case have a nice explicit form (see [C-Hisil, AsiaCrypt'17] and [Renes, PQCrypto'18]). **Thus, this is a call for fast (ℓ, ℓ) -isogenies on fast Kummers!**
- Gut feeling is that there's a better way to write down supersingular Kummers, and their arithmetic. **Thus, this is a call for smart geometers!**

Cheers!



<https://eprint.iacr.org/2018/850.pdf>

<https://www.microsoft.com/en-us/download/details.aspx?id=57309>